



TeRM: Extending RDMA-Attached Memory with SSD

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RDMA-based Storage System

- **RDMA catalyzes in-memory storage systems**
 - File systems, key-value stores, transactional databases, ...

Assise [OSDI'20]

Pilaf [ATC'13]

FaRM [NSDI'14]

Aurogon [FAST'22]

Octopus [ATC'17]

Cell [ATC'16]

Sherman [SIGMOD'22]

TH-DPMS [TOS'20]

Orion [FAST'19]

XStore [OSDI'20]

Rowan [OSDI'23]

FileMR [NSDI'20]

RACE [ATC'21]

FORD [FAST'22]

DrTM+H [OSDI'18]

ROLEX [FAST'23]

FUSEE [FAST'23]

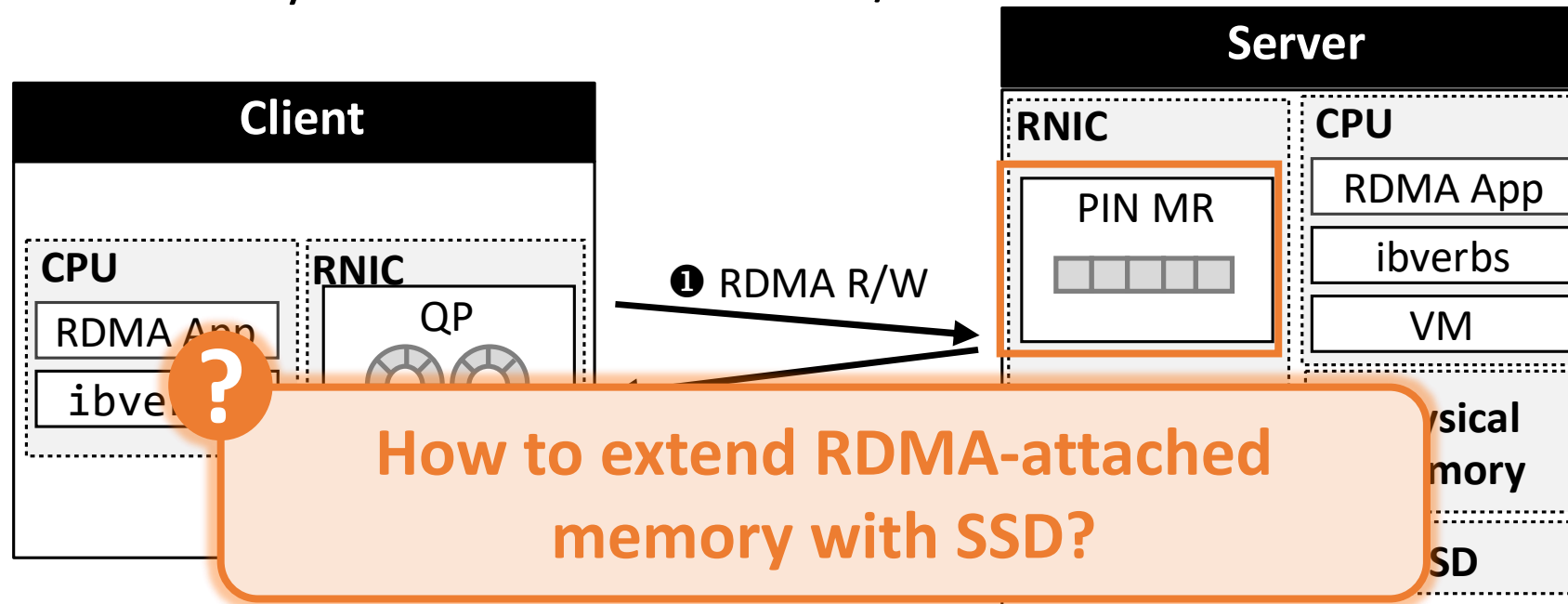
RDMA-attached Memory

- **Server**

- Expose virtual memory via **RDMA MR (RDMA-attached Memory)**
- RNIC accesses the virtual memory via DMA, bypassing the CPU
- Pin pages in the physical memory; build the RNIC page table

- **Client**

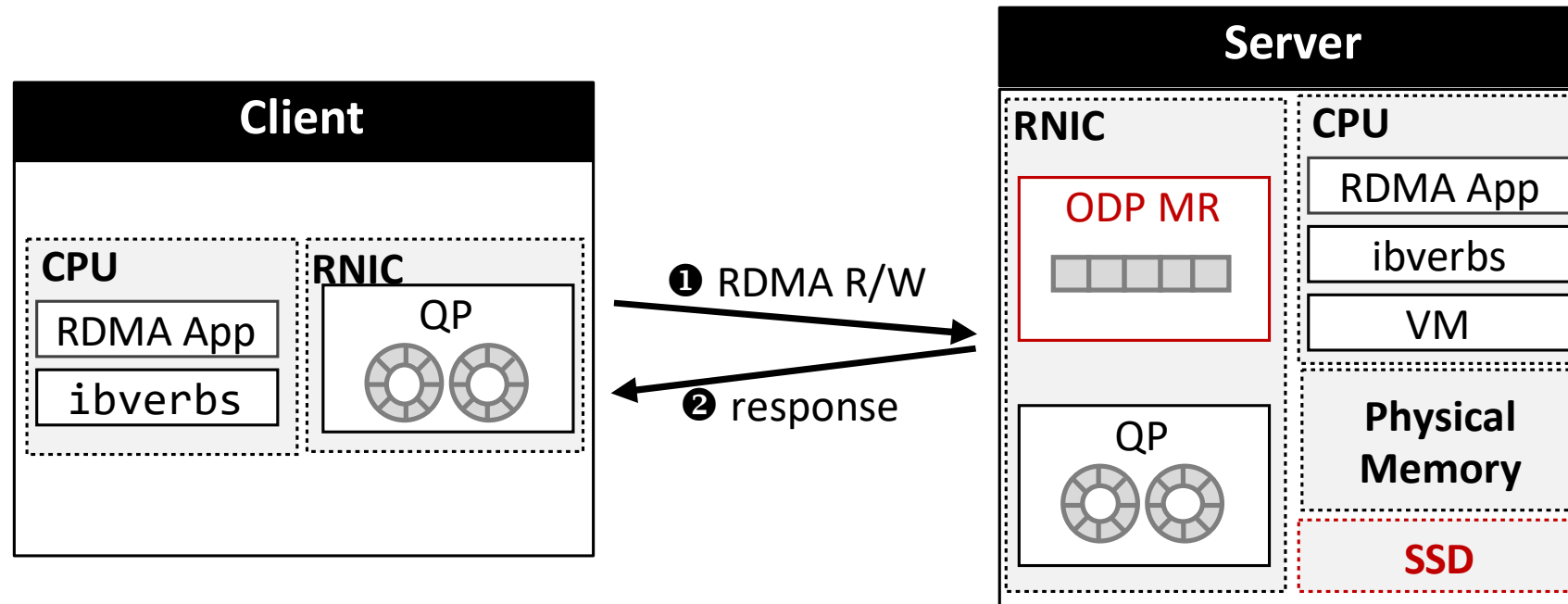
- Access the MR by one-sided RDMA READ/WRITE



ODP MR

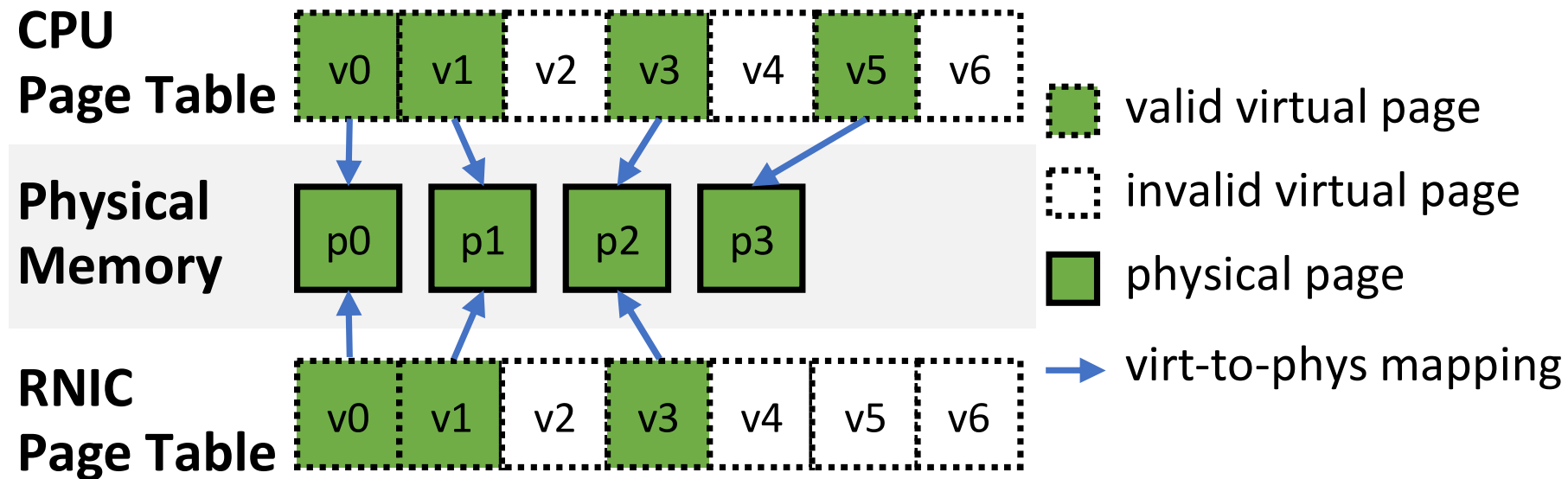
• On-demand Paging MR

- Hardware solution by Mellanox [ASPLOS'17]
- mmap an SSD and register as an ODP MR
- The client submits normal RDMA READ/WRITE



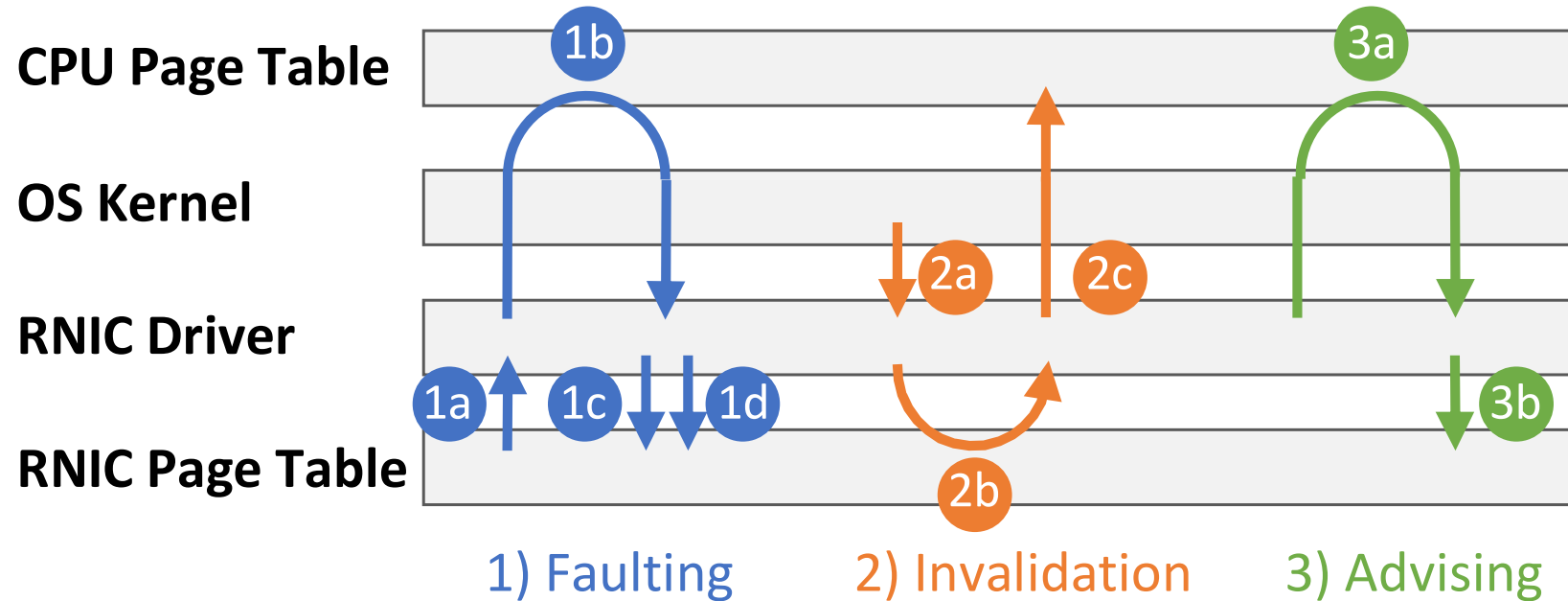
ODP MR

- Not all pages are mapped
- Trigger an **RNIC page fault** when accessing an invalid virtual page



ODP MR

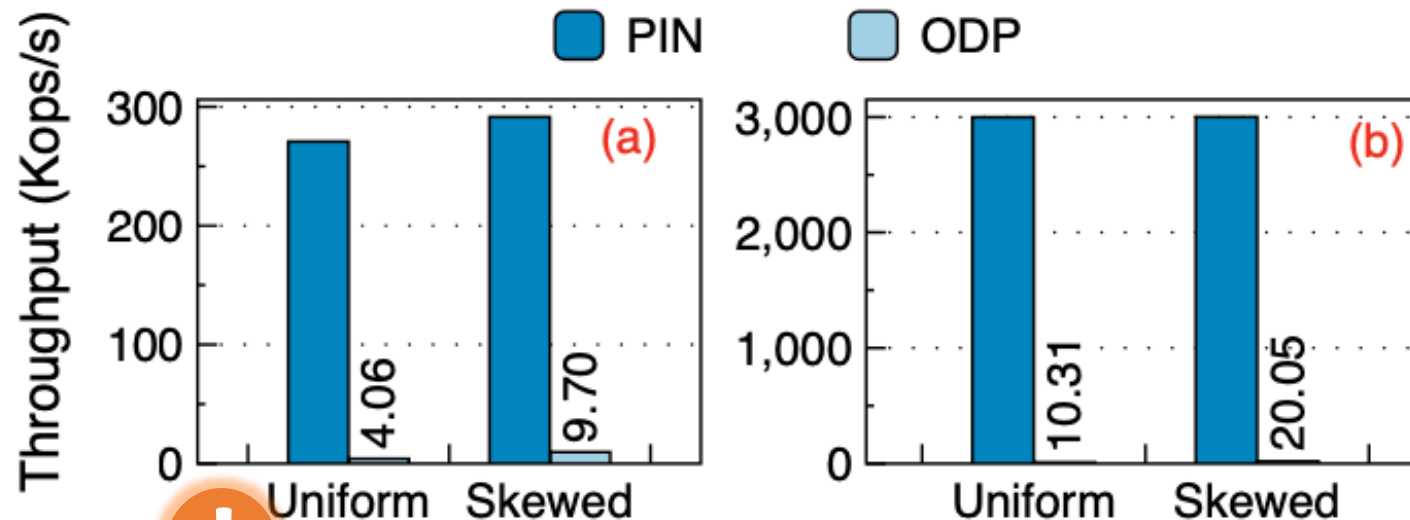
- **Synchronizing between CPU and RNIC page tables**
 - Three flows: faulting, invalidation, advising



ODP MR is not the silver bullet

- **Read 4KB performance**

- 64GB virtual memory , 32GB physical memory
- mmap() Intel Optane P5800X SSD
- (a) 1 client thread
- (b) 64 client threads

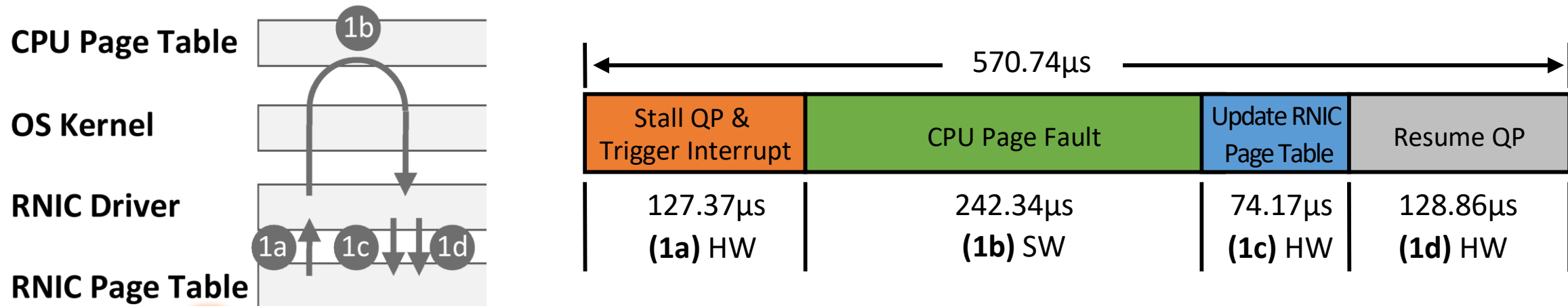


66.64x – 290.76x slowdown!

ODP MR is not the silver bullet

- **Two sources of overhead**

- A normal read consumes 4μs
- Hardware: stall & resume QP, trigger interrupt, update RNIC page table
- Software: CPU page fault



1. Onload exception handling from HW to SW.
2. Eliminate CPU page faults from the critical path.

TeRM overview

- **CPU VM**

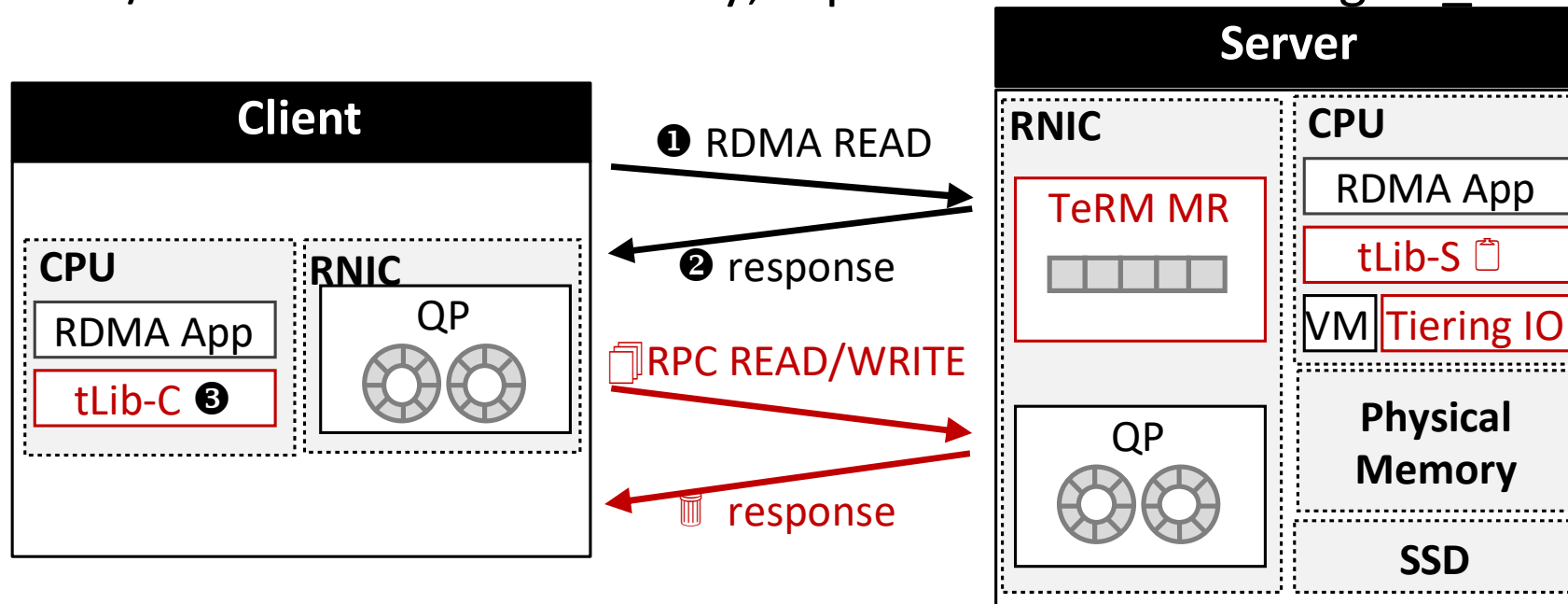
- mmap; Serves local access (load/store) from the server-side application.

- **TeRM MR**

- Serves remote access (memory read/write) from the client-side application.

- **tLib-S/tLib-C**

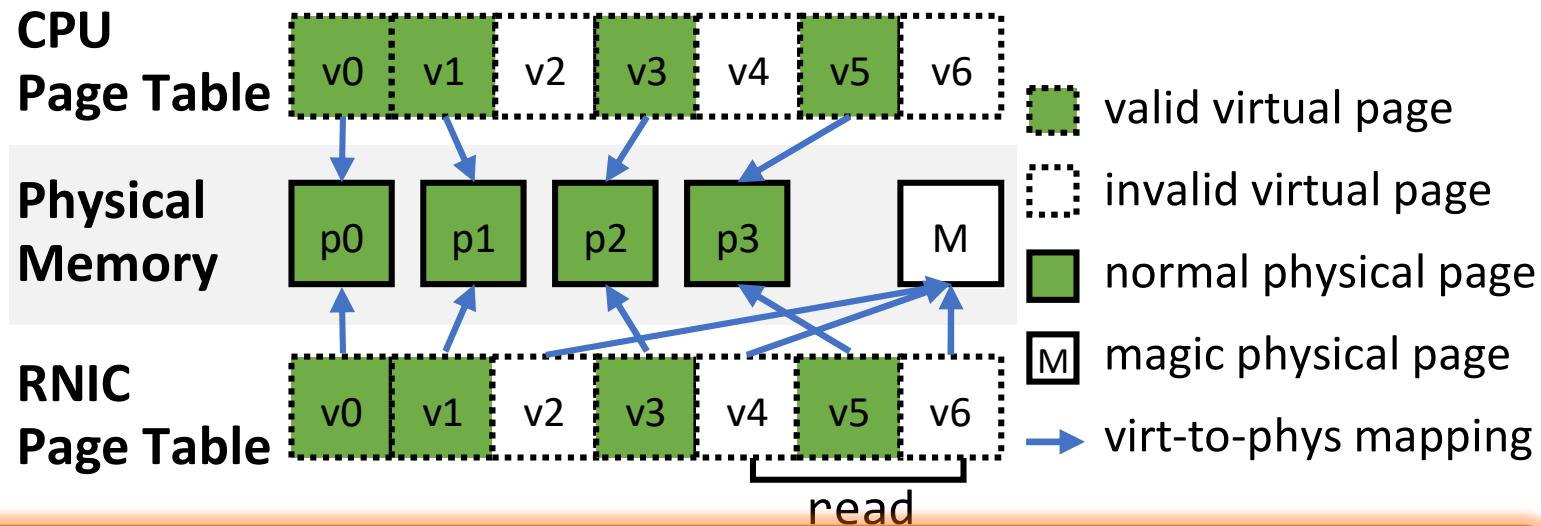
- Server-side/client-side shared library; replaces libibverbs using LD_PRELOAD



TeRM MR

- **Magic physical page**

- Invalid virtual pages are mapped to this one.
- Filled with magic pattern.



RDMA READ on invalid virtual pages returns with magic pattern.

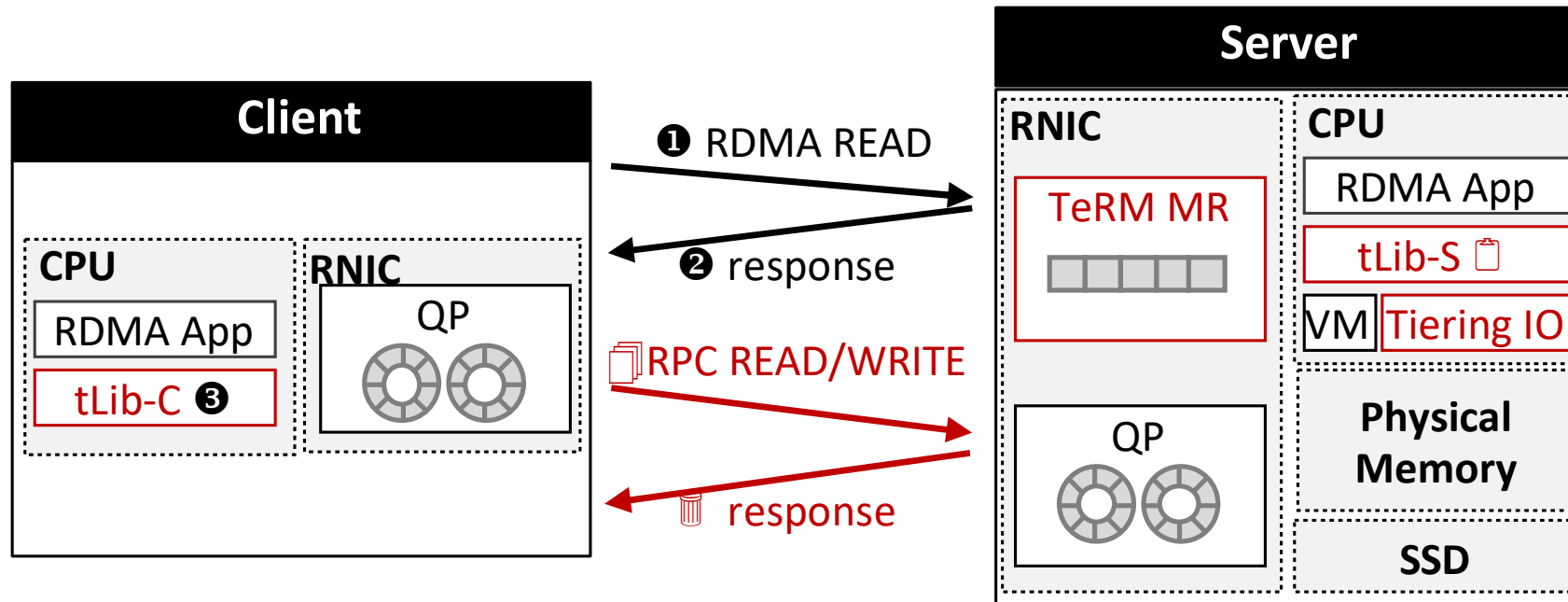
Read workflow

- **RDMA READ first**

- ❶ submit an RDMA READ request
- ❷ receive the response
- ❸ check whether the data contains magic pattern

If no magic pattern is found, the read request completes.

Otherwise, ...

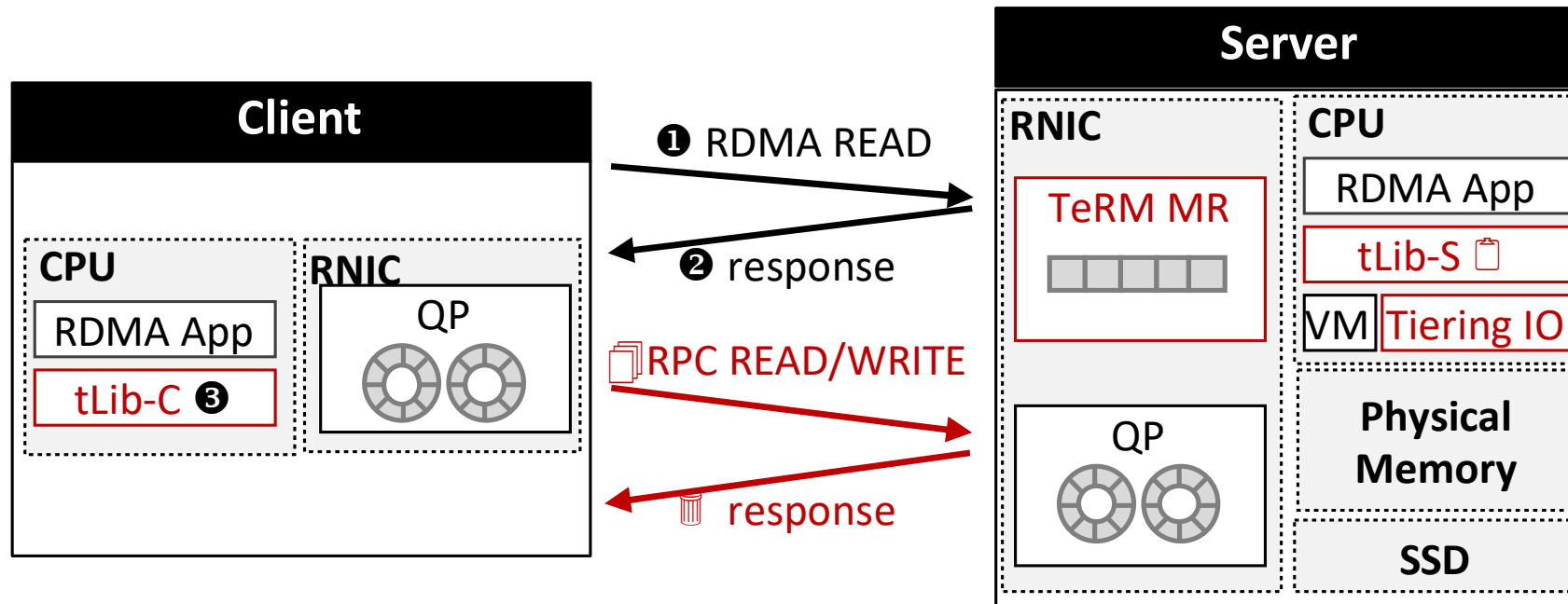


Read workflow

- **RPC READ if necessary**

- ① submit an RPC READ request
- ② tLib-S reads data
- ③ tLib-C receives data and completes the read

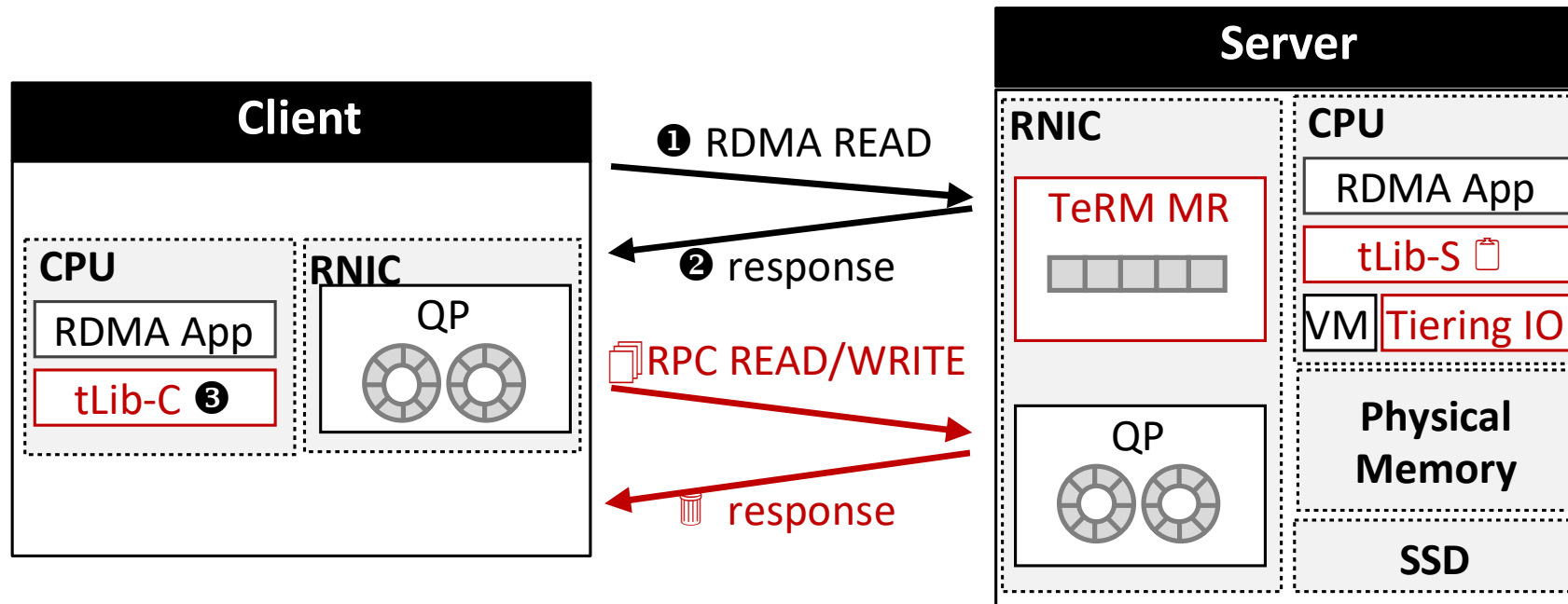
“principle 1: onload exception handling from HW to SW”



Write workflow

- **RPC WRITE for all**

- ① submit an RPC WRITE request
- ② tLib-S writes data
- ③ tLib-C receives notification and completes the write



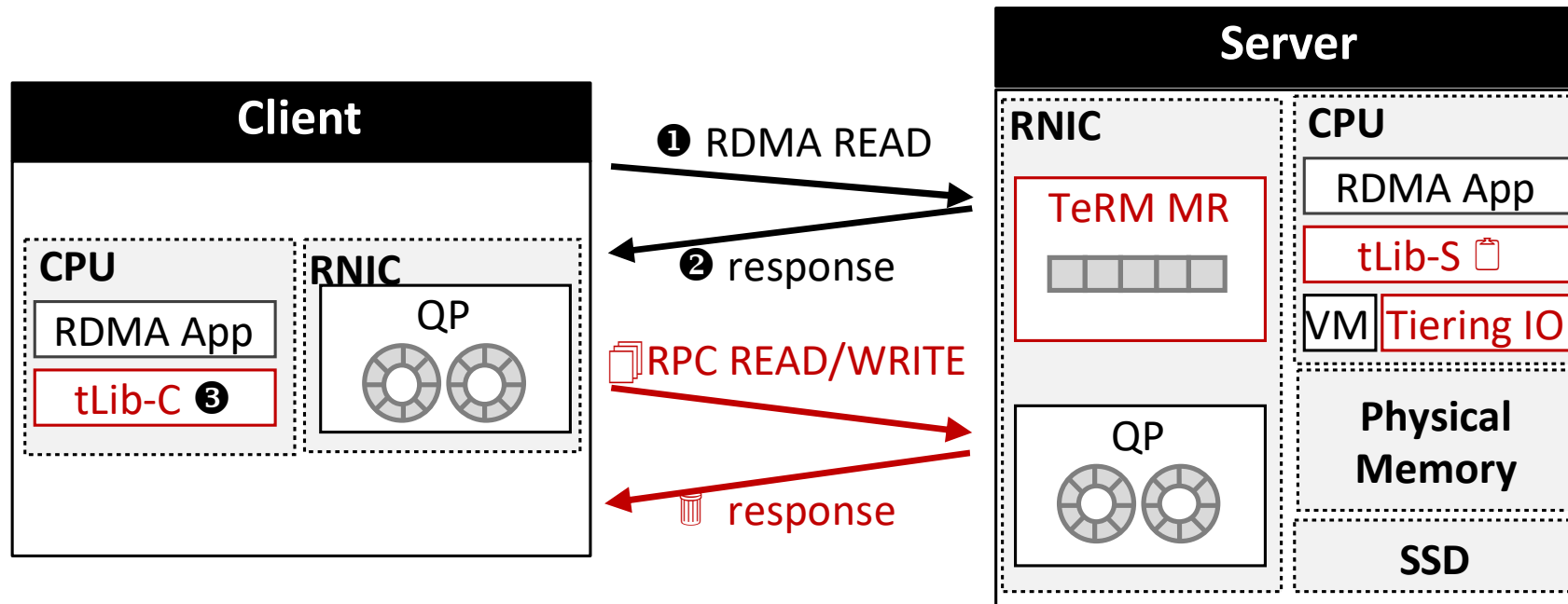
How can RPC access data efficiently?

- **Load/store the CPU VM?**

- Still triggers CPU page faults!

- **Convert memory load/store to file I/O**

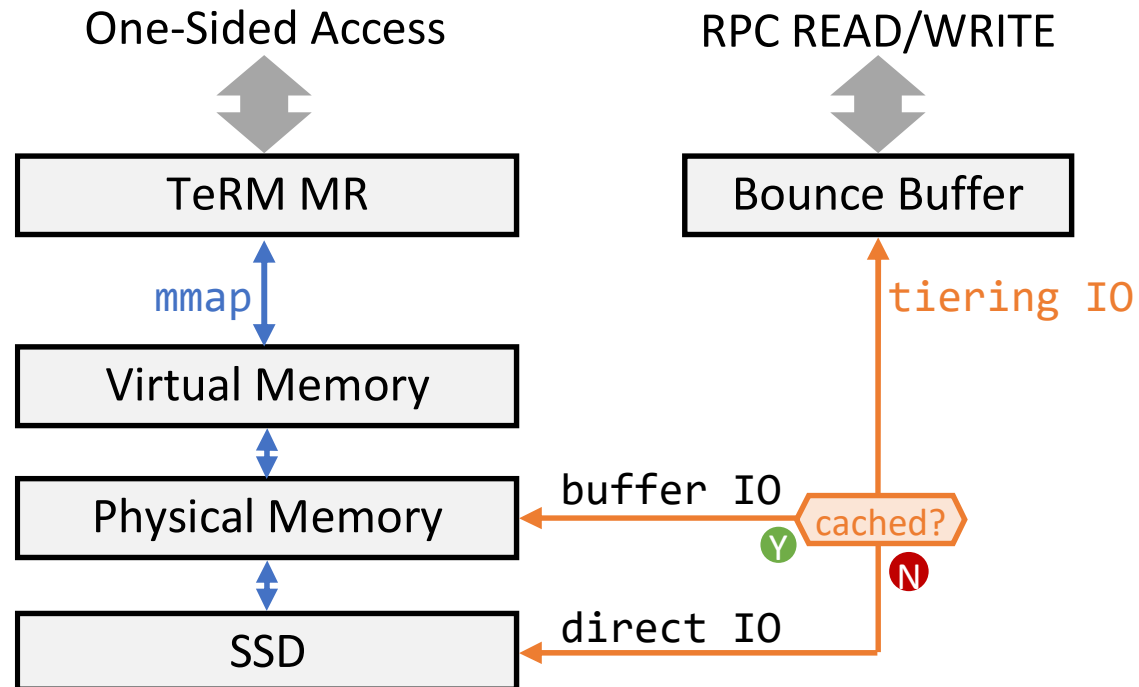
- Read/write the SSD
- “Principle 2: eliminate CPU page faults from the critical path”



How can RPC access data efficiently?

- **Convert memory load/store to file I/O**

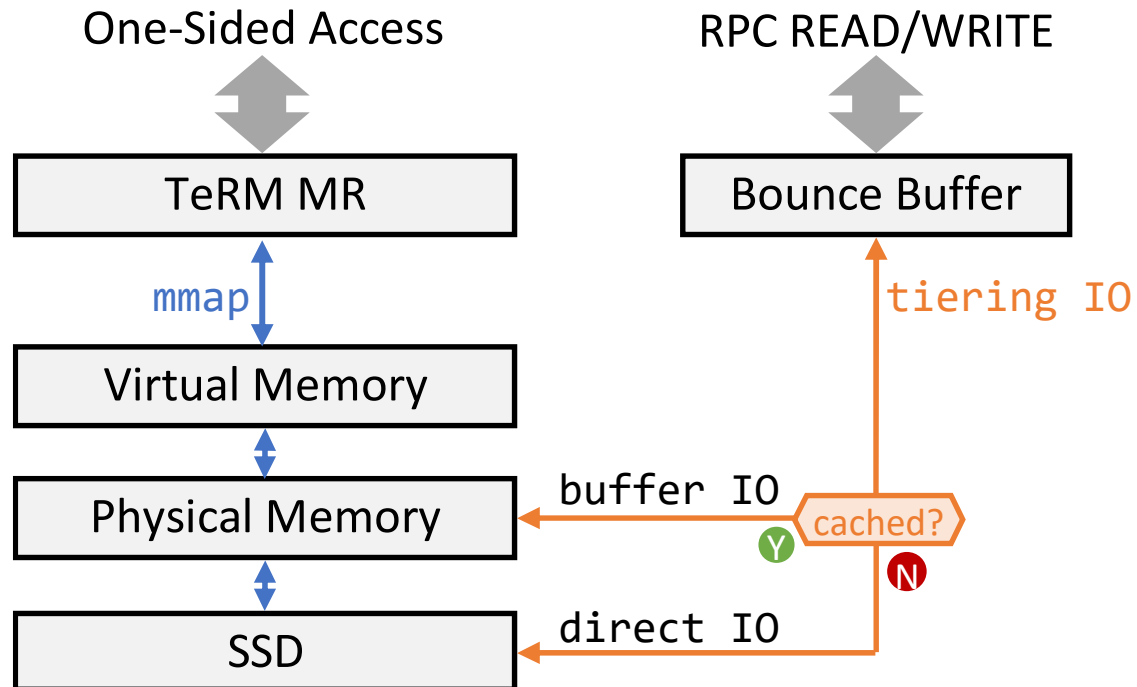
- SSD LBA range: [slba, slba + length)
- Virtual address range: [saddr, saddr + length)
- lba = addr – saddr + slba



Tiering IO

- **Read/write data via two interfaces**

- Check the page cache
- Buffer IO for cached data, using page cache
- Direct IO for uncached data, bypassing page cache



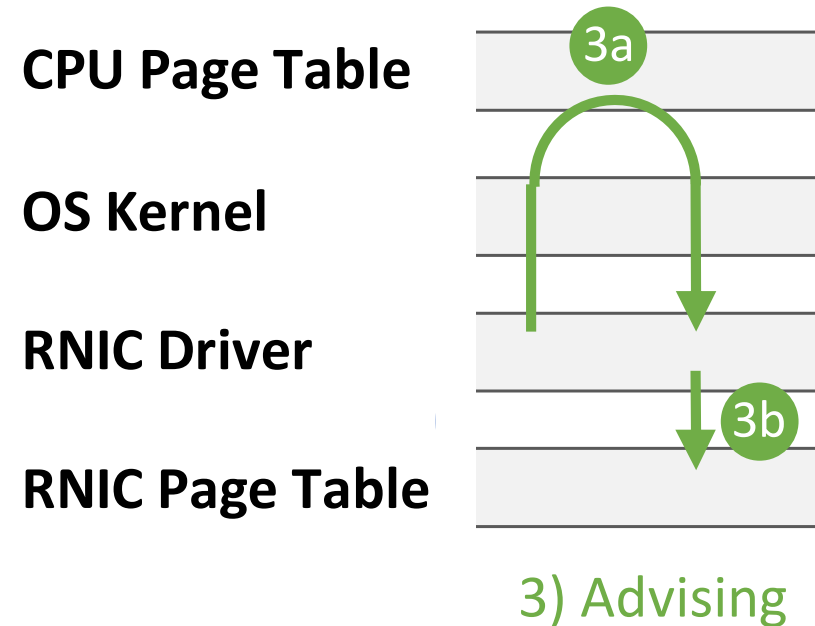
Promoting Hotspots

- **Client-side**

- Count accesses on each unit

- **Server-side**

- Aggregate counters from all clients
- Find most-accessed units as hotspots
- Promote via `ibv_advise_mr()`



Evaluation

- **Testbed**

- RDMA Cluster: server machine * 1, client machine * 2
- SSD: Intel Optane P5800X 400GB
- RNIC: ConnectX-5 100Gbps
- Switch: IB 100Gbps

- **Settings**

- Virtual memory: 64GB, physical memory: 32GB
- 64 Client threads, 16 server threads

Evaluation

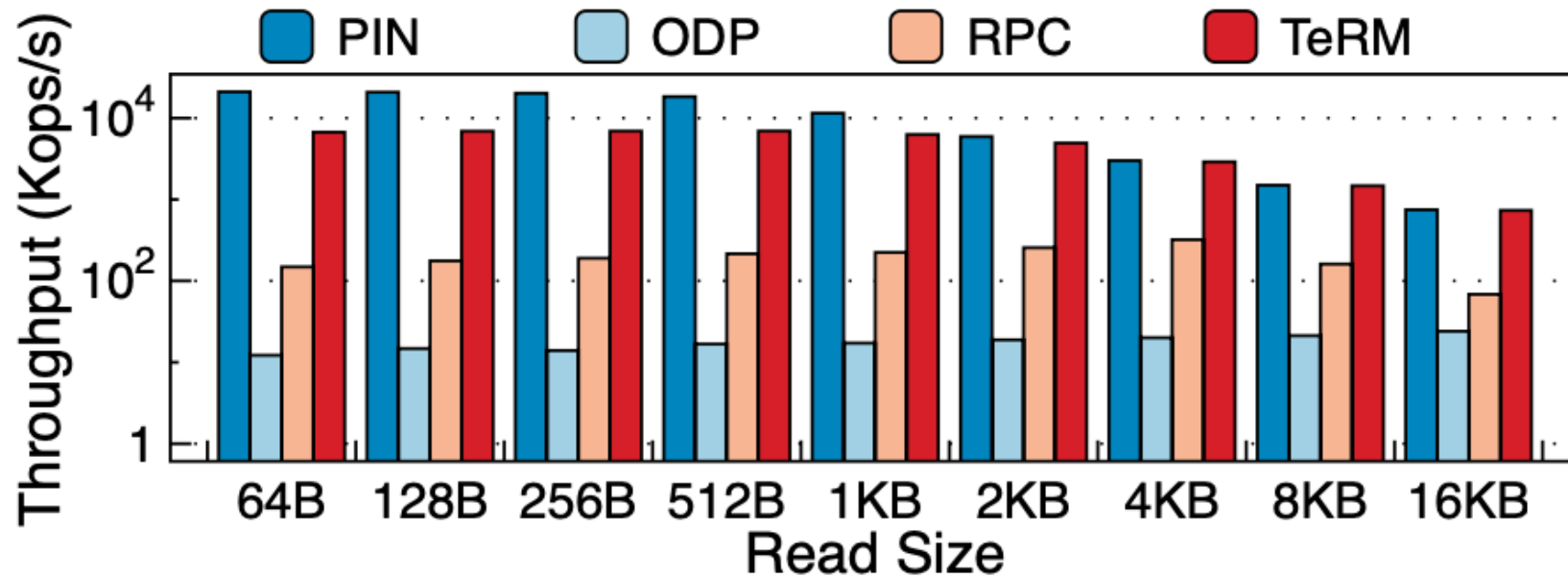
- **Comparing Targets**

- **PIN:** ideal upper bound, all pages in the physical memory
- **ODP:** hardware solution, ODP MR
- **RPC:** software solution, all requests via RPC, access data via memcpy
- **TeRM:** our solution.

Evaluation: Overall Performance

- **Read**

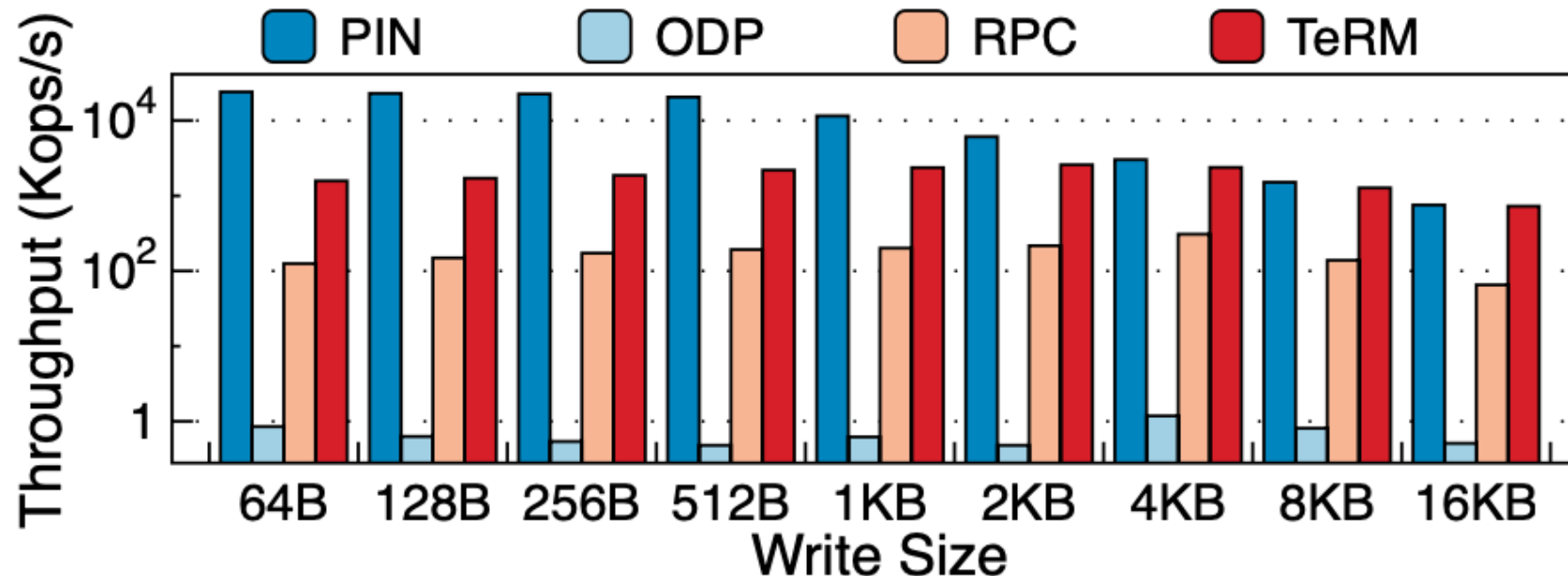
- **vs. ODP:** 30.46x – 549.63x
- **vs. RPC:** 9.05x – 45.19x
- **vs. PIN:** 37.79% – 96.71%



Evaluation: Overall Performance

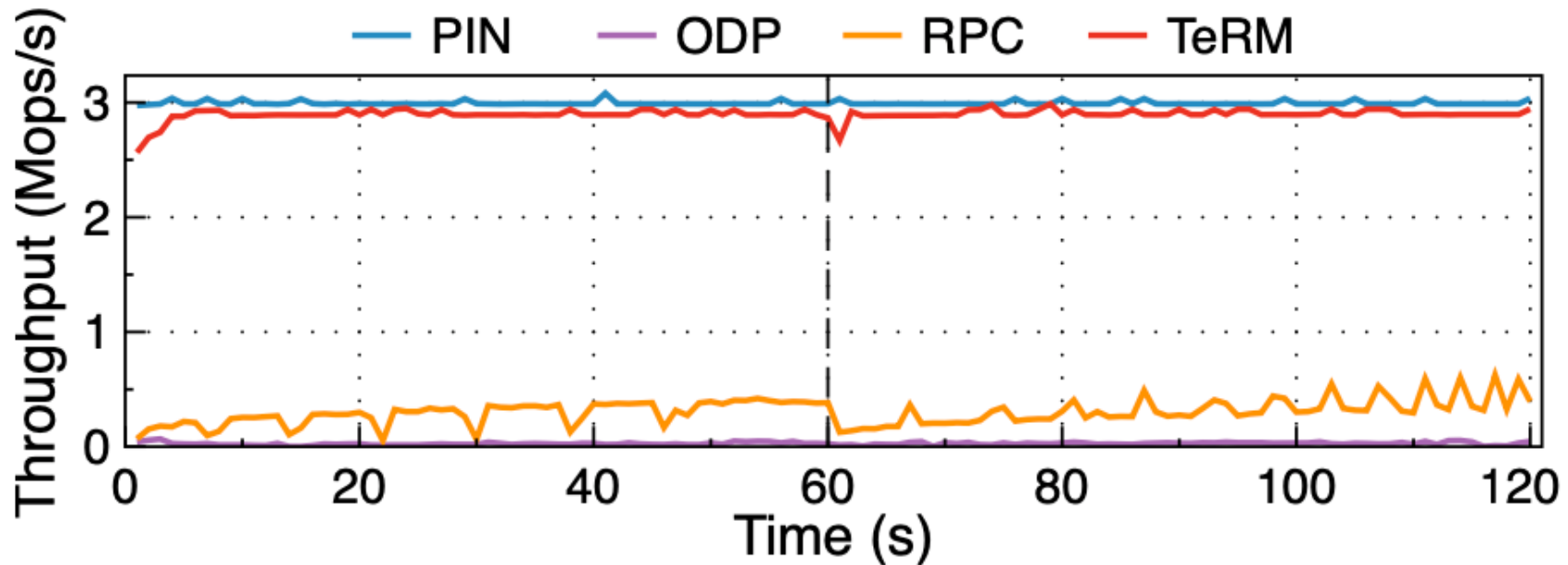
- Write

- vs. ODP: ~ 1000x (ODP write is very unstable and jitters sharply)
- vs. RPC: 7.73x – 12.60x
- vs. PIN: 6.55% – 96.32%



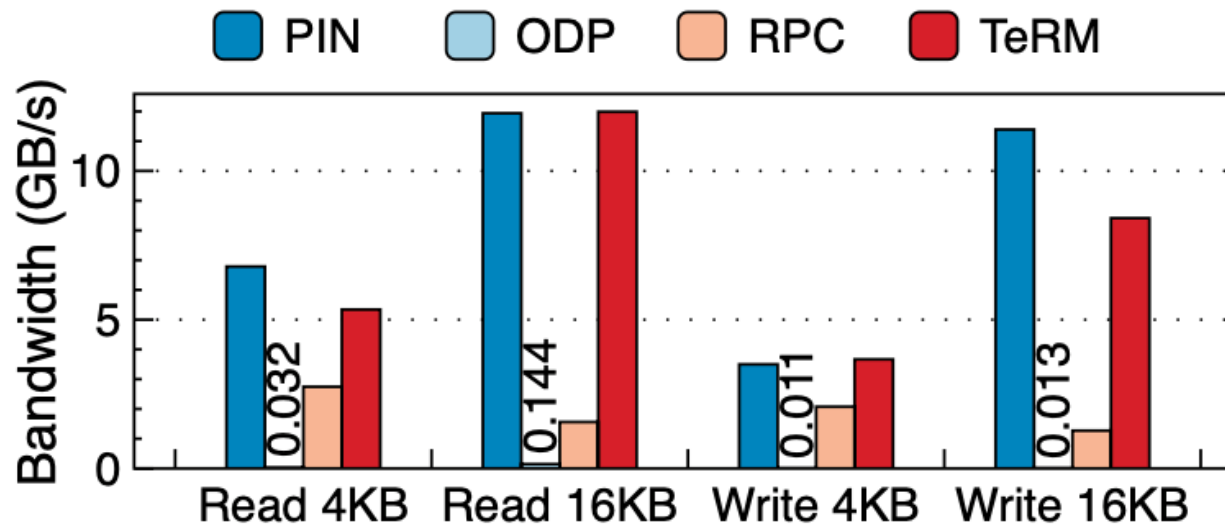
Evaluation: Dynamic Workloads

- **Change hotspots at the 60th second**
 - **Performs stably:** drops by only **6.82%**
 - **Promoting fast:** returns to the peak in **1 second**



Evaluation: RDMA-based storage system

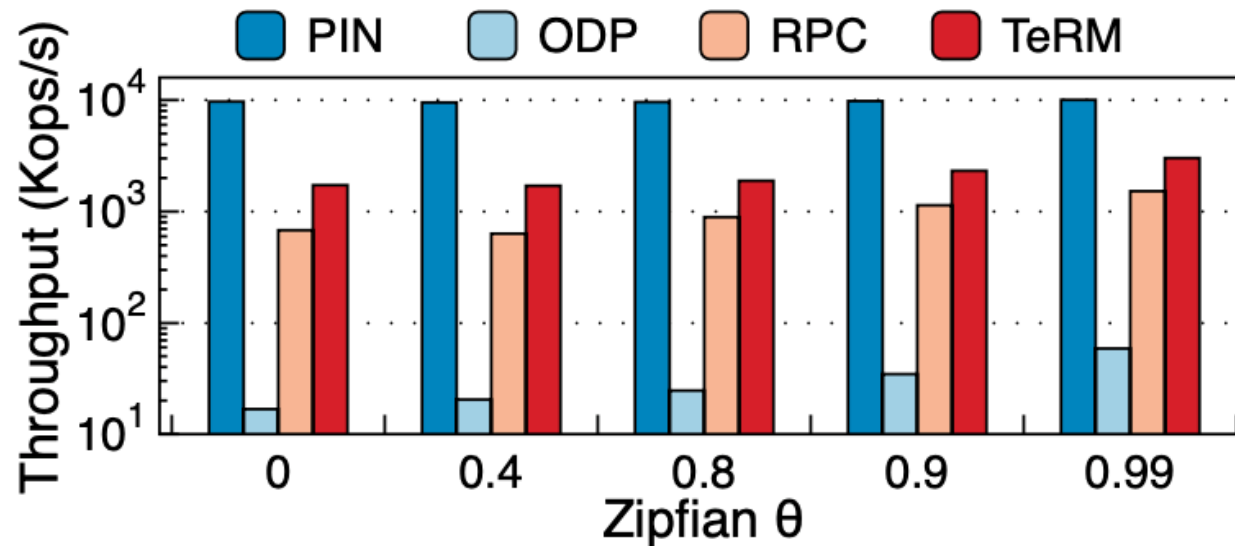
- **Octopus: A File System [ATC'17]**
 - **Workloads:** read/write the file
 - **Results:** up to **642.23x ODP**, **7.68x RPC**



Evaluation: RDMA-based storage system

- **XStore: A Key-Value System [OSDI'20]**

- **Workloads:** YCSB-C, read 8B keys and 128B values
- **Results:** Up to **102.97x ODP**, **2.69x RPC**



Conclusion

- TeRM proposes an efficient approach to **extending RDMA-attached memory with SSD**.
- TeRM **onloads exception handling** from hardware to software and eliminates RNIC & CPU page faults on the critical path.
- TeRM implements **a userspace shared library** to replace libibverbs and run **unmodified RDMA applications** transparently.
- TeRM outperforms the **hardware-only** ODP MR by up to **642.23x**, and the **software-only** RPC approach by up to **7.68x**.

Thanks! Q&A



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<https://github.com/thustorage/TeRM>